Digital Logic and Representation

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Techniques of Geoscientific Experimentation

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The way we think of and store numbers and do logic is much different than how computers do



In digital logic there are two states



Image: <u>bobvila.com</u>

In reality we have to define ranges of voltages that represent high and low states



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In reality we have to define ranges of voltages that represent high and low states $5V - V_{cc}$



Often we find the need to shift between these logic levels, which can be accomplished with a variety of techniques



To choose a resistance value for R1, consider the following equation:

 $Minimum \ Resistance \ R1 = \frac{|Signal \ Maximum \ \Delta V|}{Maximum \ allowed \ ESD \ diode \ current}$

Images: <u>sparkfun.com</u>, <u>parallax.com</u>

Pins are often marked as active low or active high



SN54HC165, SN74HC165

www.ti.com

SCLS116G - DECEMBER 1982 - REVISED AUGUST 2013

8-BIT PARALLEL-LOAD SHIFT REGISTERS

Check for Samples: SN54HC165, SN74HC165



There are a few digital logic operations that make up all of how we do computing with binary information



We show the way each operation works with a truth table



NOT gates invert the input



AND gates are only true if both inputs are true

Written AB A • B



OR gates are only true if either or both inputs are true





XOR gates are only true if either input is true





NOR gates are only true if neither input is true



 $\overline{(A + B)}$ (A + B)'

Written

Image: sparkfun.com

NAND gates are true unless both inputs are true

We can connect these logic gates together to perform calculations and other functions in a combinatorial logic circuit

Image: sparkfun.com

Activity: Fill out the truth table and write an expression for the output of this combinatorial circuit

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We can also make a sequential circuit that uses memory and (generally) a clock signal

D type flip-flops latch the input to the outputs on a clock

T flip-flops toggle the output if T and do nothing if not T

JK flip-flops can set, clear, or toggle their outputs

We can perform logic operations in software as well

1 byte a = b01010101; 2 byte b = b10101010; 3 byte c; 4 5 c = a & b; // bitwise AND-ing of a and b; the result is b000000000 6 c = a | b; // bitwise OR-ing of a and b; the result is b11111111 7 c = a ^ b; // bitwise XOR-ing of a and b; the result is b11111111 8 c = ~a; // bitwise complement of a; the result is b10101010

We use such operations to manipulate data when working with registers for example

```
1 byte a = b01010101;
2 byte b = b10101010;
3 byte c;
4
5 c = b000001111 & a; // clear the high nibble of a, but leave the low nibble alone.
6 // the result is b00000101.
7 c = b11110000 | a; // set the high nibble of a, but leave the low nibble alone.
8 // the result is b11110101.
9 c = b11110000 ^ a; // toggle all the bits in the high nibble of a.
10 // the result is b10100101.
```

We also often bit shift values to "roll" them

1 byte d = b11010110; 2 byte e = d>>2; // right-shift d by two positions; e = b00110101 3 e = e<<3; // left-shift e by three positions; e = b10101000</pre>

A 7	A 6	A 5	A 4	A 3	A 2	A 1	A ₀
128	64	32	16	8	4	2	1

A 7	A ₆	A 5	A 4	A 3	A ₂	A 1	A ₀
128	64	32	16	8	4	2	1
0	1	0	1	1	0	1	1

A 7	A 6	A 5	A ₄	A ₃	A 2	A 1	A ₀
128	64	32	16	8	4	2	1
0	1	0	1	1	0	1	1
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A 7	A ₆	A 5	A 4	A 3	A ₂	A 1	A ₀
128	64	32	16	8	4	2	1
0	1	0	1	1	0	1	1
0	64	0	16	8	0	2	1
							91

The most significant bit can be first or last, we just have to agree and know what was done

Image: users.cis.fiu.edu

Endianness gets its name from Swift's Gulliver's Travels

BIG ENDIAN - The way people always broke their eggs in the Lilliput land

LITTLE ENDIAN - The way the king then ordered the people to break their eggs

We've explored base 10 and base 2, but what if we want more than 0-9? Base 16!

0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F

Counting is a bit strange since we're used to base 10

Decimal	Hexadecimal	•••	Decimal	Hexadecimal
0	0		8	8
1	1		9	9
2	2		10	Α
3	3		11	В
4	4		12	С
5	5		13	D
6	6		14	E
7	7		15	F

Decimal	Hexadecimal	 Decimal	Hexadecimal
16	10	24	18
17	11	25	19
18	12	26	1A
19	13	27	1B
20	14	28	1C
21	15	29	1D
22	16	30	1E
23	17	31	1F

Let's convert the decimal number 48879 to hex

$$48879/16 = 3054 \text{ R } 15 > F$$

$$3054/16 = 190 \text{ R } 14 > EF$$

$$190/16 = 11 \text{ R } 14 > EEF$$

$$11/16 = 0 \text{ R } 11 > BEEF$$

Converting to binary is a powers of 16 problem

A ₇	A ₆	A 5	A ₄	A ₃	A ₂	A ₁	A ₀
16 ⁷	16 ⁶	16 ⁵	16 ⁴	16 ³	16 ²	16 ¹	16 ⁰
268435456	16777216	1048576	65536	4096	256	16	1

Converting from binary is done by grouping into bunches of 4

Binary: 1011111011101111 Binary: 1011 | 1110 | 1110 | 1111 Decimal: 11 | 14 | 14 | 15 HEX: B | E | E | F

You'll see all of these formats written in a variety of ways

Hex OxBEEF #FF7454 %20 \x1B &#BD **OhBEEF** BEEF₁₆ BEEFhex

Two's complement lets us represent negative numbers

Converting to two's compliment is "simple"

Write out the number in binary Invert all of the digits Add one to the result

Write the number -42

- 1. Write out the number in binary
- 2. Invert all of the digits
- 3. Add one to the result

42 = 2 + 8 + 32 0010_1010

Write the number -42

- 1. Write out the number in binary
- 2. Invert all of the digits
- 3. Add one to the result

0010_1010
 1101_0101

Write the number -42

- 1. Write out the number in binary
- 2. Invert all of the digits
- 3. Add one to the result

1101_0101 1101_0110

Floating point is expensive, but useful

Single-Precision

sign * 2^{exponent} * mantissa

In this example:

$$\begin{array}{l} \operatorname{sign} = b_{31} = 0 \\ \bullet \ (-1)^{\operatorname{sign}} = (-1)^0 = +1 \in \{-1, +1\} \\ \bullet \ e = b_{30}b_{29} \dots b_{23} = \sum_{i=0}^7 b_{23+i}2^{+i} = 124 \in \{1, \dots, (2^8 - 1) - 1\} = \{1, \dots, 254\} \\ \bullet \ 2^{(e-127)} = 2^{124-127} = 2^{-3} \in \{2^{-126}, \dots, 2^{127}\} \\ \bullet \ 1.b_{22}b_{21} \dots b_0 = 1 + \sum_{i=1}^{23} b_{23-i}2^{-i} = 1 + 1 \cdot 2^{-2} = 1.25 \in \{1, 1 + 2^{-23}, \dots, 2 - 2^{-23}\} \subset [1; 2 - 2^{-23}] \subset [1; 2) \end{array}$$

thus:

• value = (+1)
$$\times 1.25 \times 2^{-3} = +0.15625$$

Image: wikipedia.com

Not all representations are exact and can accumulate error

Decimal Number: 0.1 Single Precision: 0x3DCCCCCD Cast to Double: 0.1000000149011612

1e6 * 0.1 (cast to double) = 100000.00149011612

Image: wikipedia.com

Assignment: Digital Representation

DUE: 9/27/16